## **Homework 5: Solutions**

**6.7.4** Applying Hall's theorem and defining  $L = \lfloor 2 \log n \rfloor$ ,

$$\mathbf{P}(\text{no p.m.}) \leq \sum_{k=L}^{n} \binom{n}{k} \binom{n}{k-1} \left(\frac{\binom{k-1}{L}}{\binom{n}{L}}\right)^{k}$$
$$\leq \sum_{k=L}^{n} \left(\frac{ne}{k}\right)^{k} \left(\frac{ne}{k-1}\right)^{k-1} \left(\frac{k-1}{n}\right)^{Lk}$$
$$\leq \sum_{k=L}^{n} \left(\left(\frac{k}{n}\right)^{L-2} \cdot e^{2} \cdot \left(\frac{k-1}{k}\right)^{L}\right)^{k}$$
$$= \sum_{k=L}^{n} u_{k}.$$

Now if  $k = (1 - \epsilon)n$  then we have  $(k/n)^{L-2} \leq e^{-\epsilon(L-2)} \leq e^{-3+o(1)}$  if  $\epsilon \geq 3/L$ . We deduce from this that

$$\sum_{k=L}^{n_0} u_k = o(1) \text{ where } n_0 = n - \frac{4n}{L}.$$

So we need to be more careful for  $k > n_0$ . If there is a set of size k on one side of the partition with at most k - 1 neighbors, then there is a set X of size  $\ell = n - k + 1$  on the other side of the partition with at most  $\ell - 1$  neighbors. We estimate this by

$$\binom{n}{\ell-1} \mathbf{P} \left( Bin\left(n, 1 - \left(1 - \frac{\ell}{n}\right)^L\right) \le \ell - 1 \right).$$
 (1)

We choose X in  $\binom{n}{\ell-1}$  ways and then  $1 - (1 - \frac{\ell}{n})^L$  lower bounds the probability that a vertex chooses a neighbor in X.

Summing over  $\ell \leq \ell_0 = n - n_0 + 1$ , this is at most

$$2\sum_{\ell=1}^{\ell_0} {\binom{n}{\ell-1}}^2 \left(1 - \left(1 - \frac{\ell}{n}\right)^L\right)^{\ell-1} \left(1 - \frac{\ell}{n}\right)^{L(n-\ell+1)}$$
(2)  
$$\leq 2\sum_{\ell=1}^{\ell_0} \left(\frac{ne}{\ell-1}\right)^{2\ell-2} \left(\frac{\ell L}{n}\right)^{\ell-1} e^{-L\ell(1+o(1))}$$
  
$$\leq 2\sum_{\ell=1}^{\ell_0} \left(\frac{ne^2\ell L}{\ell-1}\right)^{\ell-1} n^{-(2-o(1))\ell}$$
  
$$= o(1).$$

**Explanation for** (2): the binomial probability in (1) is dominated by the  $\ell - 1$  term, leading to the factor 2.

- **6.7.9** Following the hint we partition [n] into 3 sets A, B, C of size n/3. The bipiartite graph H induced by A, B is distributed as  $G_{n/3,n/3,p}$  and since  $\frac{n}{3}p \gg \log \frac{n}{3}$  this graph has a perfect matching w.h.p. Fix a perfect matching M of H and define another random bipartite graph K with **vertices** M, C and an edge (e, x) for each  $e = \{u, v\} \in M, x \in C$  such that the edges  $\{x, u\}, \{x, v\}$  both exist. The random graph K is distributed as  $G_{n/3,n/3,p^2}$  and since  $\frac{n}{3}p^2 \gg \log \frac{n}{3}$  this graph has a perfect matching w.h.p. This perfect matching corresponds to n/3 vertex disjoint triangles.
- 6.7.10 Arguing as in the proof of Theorem 6.1 we see that

$$\mathbf{E}Y \le 2\sum_{k=2}^{n/2} \binom{n}{k} \binom{n}{k-1} \binom{k(k-1)}{2k-2} p^{2k-2} (1-p)^{k(3n/4-k)}$$

The only change here is that we can only guarantee that S has at least k(3n/4 - k) neighbors not in T. Continuing,

$$\begin{split} \mathbf{E}Y &\leq 2\sum_{k=2}^{n/2} \left(\frac{ne}{k}\right)^k \left(\frac{ne}{k-1}\right)^{k-1} \left(\frac{Kke\log n}{2n}\right)^{2k-2} n^{-Kk(3/4-k/n)} \\ &\leq \frac{n^2}{\log^2 n} \sum_{k=2}^{n/2} \left(\frac{ne}{k} \cdot \frac{ne}{k-1} \cdot \left(\frac{Kke\log n}{2n}\right)^2 \cdot n^{-K/2}\right)^k \\ &= o(1), \end{split}$$

if  $K \geq 4$ .