Operations Research II Final Project

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1 Introduction

1.1 Motivation

Carnegie Mellon has a rigorous Mathematics degree program regardless of the classes a student takes over the course of their 4 years here. Students, every semester, have to make decisions on what classes to take while considering interests, major requirements, and difficulty of classes. All else aside, the difficulty of classes has a impact on the student experience and ability to participate in activities outside of the classroom. The goal of this project is to create a schedule that minimizes the Faculty Course Evaluation (FCE) hours, which are the average number of hours spent by students on all course-related activities for a given class. Specifically, we aim to minimize the maximum number of FCE hours in any semester.

FCE hours provide a useful way to measure how difficult and time consuming a class actually is based on actual student feedback. Thus, we thought it would be a sufficient measure to minimize. Our 8-semester schedule can be used as a guide for future students to plan out a manageable schedule for their time at CMU.

1.2 Previous Projects

One of our goals is to improve and expand upon previous projects of a similar topic. One project, from Fall '19, has provided a solution, except it has some assumptions that we have made stronger. For instance, in their paper, they note that they assume any class can be taken any semester - which is not the case and provides solutions that are not reasonable. In addition, their solution is solely math classes. Our project, instead, focuses on the core curriculum of the Mathematical Science (Statistics and Operations Research concentration)

and provides a more thorough planning guide which includes statistics, business, and computer science courses as options for the depth electives. Furthermore, their provided solution includes taking 700 level courses as a first-year student. This is highly unlikely to do, and our solution will provide a more practical solution for the general student. Therefore, we have excluded graduate level courses and included solely undergraduate level classes.

Another project, from Fall '18, provides a solution, but it only contains classes for 4 semesters. Also, this schedule has prerequisite conflicts such as taking 21-270 and 21-370 in the same semester. This is also an issue as 21-370 is only offered in the Fall. Our project will focus on creating a schedule for the full 8 semesters, as well as ensuring that all prerequisite requirements are met to provide a more accurate solution to the problem. The solution provided in the past project includes taking up to 6 math classes a semester, which is not manageable for the average student. We will focus on balancing the quantity of math courses by including other major requirements and Gen-Ed courses. Overall, improving the faults in these two projects were a focus of our project.

2 Formulating ILP

2.1 Decision Variables

We first want to create decision variables. We are trying to figure out the optimal schedule, so the decision variables will center around which courses to take and when. Specifically, we have the binary variable $x_{i,j}$ to represent whether course i is taken during semester j or not. We will let N denote the number of courses and S denote the number of semesters, and thus $x_{i,j}$ is defined for $i \in |N|$ and $j \in |S|$.

2.2 Objective Function

Our goal is for the student to not have any semesters that are too busy. To that end, our objective is to minimize the maximum FCE in any given semester. Thus, we let f_i represent the FCE of course i for $i \in [N]$. Then the objective function can be written as

$$
\min \quad \max_{j \in S} \left\{ \sum_{i=1}^N f_i \cdot x_{i,j} \right\}.
$$

We run into a problem however. Our goal is to have an integer linear program, but the max function isn't directly linear, and would therefore be hard to encode into the software we will use. Hence, we create one more decision variable a which will represent the max FCE of any semester. Then our objective is to minimize a , but we must also add constraints ensuring that a is at least the total FCE of each semester:

$$
a \ge \sum_{i=1}^N f_i \cdot x_{i,j} \quad \forall j \in [S].
$$

Then in a feasible solution, a will not necessarily represent the maximum FCE of a semester (it could be higher than the maximum), but in any optimal solution, it will.

2.3 Constraints

We have quite a few constraints to add. First of all, any particular class can not be taken more than once. So

$$
\sum_{j=1}^{S} x_{i,j} \le 1 \quad \forall i \in [N].
$$

Moreover, there is a maximum number of units that a student is allowed to take per semester before their schedule is considered an overload. We will call this variable U_o (for the average Mellon College of Science student, $U_o = 54$). Then to ensure that no semester is more than U_o units, we let u_i represent the number of units that course i is worth. We obtain the constraint:

$$
\sum_{i=1}^{N} u_i \cdot x_{i,j} \le U_o \quad \forall j \in [S].
$$

Likewise, in order to graduate, there is a minimum number of units a student must take during their time at college, which we will call U_q (for CMU students, $U_q = 360$ units). To ensure that at least this many units are taken, we have the constraint:

$$
\sum_{j=1}^{S} \sum_{i=1}^{N} u_i \cdot x_{i,j} \ge U_g
$$

We must then deal with prerequisites. Many courses have multiple courses that could fulfill the prerequisites. For example, the probability course 21-325 has a prerequisite of 21-268 or 21-259 or 21-269 or 21-256 (which are all 3D calculus courses). We dealt with this by creating groups of classes, labeled $G_1, G_2, \ldots, G_\gamma$, where γ is the total number of groups. Then we have the set $P = \{(i, g) \in [N] \times [\gamma] : \text{course } i \text{ has group } G_g \text{ as a prerequisite}\}.$ So in our example, we would put the four calculus courses above in one group (say G_1), and then there would be the ordered pair $(21325, 1) \in P$ representing the fact that a course from group G_1 must be taken as a prereq for the course 21-325. Importantly, only one course

from the group needs to be taken, not multiple or all of them, and this will always be the case. We therefore end up with the constraint:

$$
\sum_{k=1}^{j-1} \sum_{h \in G_g} x_{h,k} \ge x_{i,j} \quad \forall j \in [S], \ \forall (i,g) \in P
$$

which dictates that if course i is taken in semester j, and if group G_g is a prerequisite for course i, then at least one course h in the group G_g must be taken in a semester before j (so notably a semester k between 1 and $j - 1$, inclusive).

We also have to encode the classes required to graduate. We define requirements R_1, R_2, \ldots, R_ρ , which are groups of required classes. Each set R_r contains all course numbers $i \in [N]$ that count towards requirement r, and in order for the requirement to be satisfied, at least one course $i \in R_r$ must be taken (importantly, not all courses have to be taken, one is enough). For example, an Operations and Statistics concentration student must take one probability course, but they can chose between 21-325, 15-259, and 36-218 to fulfill this requirement. So we would put all three of these courses in one requirement R_r . Then to make sure all requirements are met, we have the constraints:

$$
\sum_{j=1}^{S} \sum_{i \in R_r} x_{i,j} \ge 1 \quad \forall r \in [\rho].
$$

There is however one requirement that can't be formulated as above: a student must take at least 45 units of depth electives, from a long list of depth electives. For this, we define the set D to be all course numbers $i \in [N]$ that count as depth electives, and then we have the constraint:

$$
\sum_{j=1}^{S} \sum_{i \in D} x_{i,j} \cdot u_i \ge 45.
$$

The final set of constraints we must consider concern the times during which courses are taken. We need to make sure that students aren't enrolled in two courses that have lectures at the same time, and we also must ensure that a student doesn't plan to take a Spring-only class in the Fall or a Fall-only class in the Spring. To do this, we divide the week into T time slots. For example, at CMU, classes start every half hour, so we divided the week into 30-minute time slots, ranging from 8-8:30am on Monday to 8:30-9pm on Friday. We also assume that any class offered in the Spring will be at the same time every Spring, and that any class offered in the Fall will be at the same time every Fall. Then we have the binary data variables $y_{i,1,t}$ and $y_{i,2,t}$ which are 1 if course i is offered during time slot t in the Fall (season 1), and Spring (season 2), respectively, and 0 otherwise. To make sure only one class is taken per time slot per semester, we add the constraints:

$$
\sum_{i=1}^{N} x_{i,j} \cdot y_{i,2-j} \cdot z_{i,1} \le 1 \quad \forall t \in [T], \ \forall j \in [S]
$$

which sum over all courses that could be offered in the given time slot and semester and make sure at most one of them are taken. In this context, we are using $\%$ as the mod operator to denote the remainder when j is divided by 2, so that $2 - j\%2$ evaluates to 1 when j is odd (since the odd semesters are in the fall) and 2 when j is even.

Finally, we ensure that no Spring-only class is taken in the fall, and vice-versa. If a course i is offered in the Fall, it will meet during at least one fall time slot, and then $\sum_{t=1}^{T} y_{i,1,t} \geq 1$, whereas if course *i* isn't offered in the Fall, then we have $\sum_{t=1}^{T} y_{i,1,t} = 0$. If course *i* isn't offered in the Fall, we want $x_{i,1}, x_{i,3}, \ldots, x_{i,2[s/2]-1}$ to all be 0, which we can encode by forcing their sum to be 0. And we know their sum will be at most 1 (we encoded this above because each course can be taken at most 1), so enforcing that the sum is less than or equal to 1 wouldn't have any consequences. Thus we add the constraints

$$
\sum_{k=1}^{\lceil S/2 \rceil} x_{i,2k-1} \le \sum_{t=1}^{T} y_{i,1,t} \quad \forall i \in [N]
$$

to make sure that course i won't be taken in the Fall if it isn't offered in the Fall. Likewise, we have the constraints

$$
\sum_{k=1}^{\lfloor S/2 \rfloor} x_{i,2k} \le \sum_{t=1}^T y_{i,2,t} \quad \forall i \in [N]
$$

for courses that aren't offered in the Spring.

2.4 Overall Variable Definitions

After all this work, we end up with the following variables. Of the variables below, only $x_{i,j}$ and a are the decision variables, all other variables are *data variables*, meaning that they are being input into the program, their values aren't being changed to find an optimal solution.

> $N =$ the total number of courses $S =$ number of semesters $x_{i,j} =$ $\int 1$, course $i \in [N]$ is taken in semester $j \in [S]$ 0, otherwise f_i = the FCE of course i $a = \max$ FCE across semesters (enforced via constraints)

 u_i = how many units course *i* is U_o = max number of units in a semester to avoid overloading U_g = number of units needed to graduate $G_g=\{i\in [N]: \text{course } i \text{ is part of the } g\text{'th group of classes}\}$ γ = The number of groups G_g $P = \{(i, g) \in [N] \times [\gamma] : \text{course } i \text{ has a course from group } G_g \text{ as a prereq}\}\$ $R_r = \{i \in [N] : \text{course } i \text{ is a course that can be taken to fulfill requirement } r\}$ ρ = The number of requirements R_r $T =$ number of time slots in a week \mathcal{L}_{1} \mathcal{F} if course $i \in [N]$ is offered during time slot $t \in [T]$ in the Fall

$$
y_{i,1,t} = \begin{cases} 1, & \text{if course } i \in [N] \text{ is offered during time slot } t \in [I] \text{ in the Fall} \\ 0, & \text{otherwise} \end{cases}
$$
\n
$$
y_{i,2,t} = \begin{cases} 1, & \text{if course } i \in [N] \text{ is offered during time slot } t \in [T] \text{ in the Spring} \\ 0, & \text{otherwise} \end{cases}
$$

2.5 Overal ILP

Thus we obtain the ILP:

minimize a subject to

$$
a \geq \sum_{i=1}^{N} f_i \cdot x_{i,j} \quad \forall j \in [S] \qquad \text{(ensure a is the max FCE in any semester)}
$$
\n
$$
\sum_{j=1}^{S} x_{i,j} \leq 1 \quad \forall i \in [N] \qquad \text{(no class is taken multiple times)}
$$
\n
$$
\sum_{i=1}^{N} u_i \cdot x_{i,j} \leq U_o \quad \forall j \in [S] \qquad \text{(no semester is a unit overload)}
$$
\n
$$
\sum_{j=1}^{S} \sum_{i=1}^{N} u_i \cdot x_{i,j} \geq U_g \qquad \text{(the required number of units to graduate are taken)}
$$
\n
$$
\sum_{k=1}^{j-1} \sum_{h \in G_g} x_{h,k} \geq x_{i,j} \quad \forall j \in [S], \ \forall (i,g) \in P \qquad \text{(prerequires are satisfied)}
$$
\n
$$
\sum_{j=1}^{S} \sum_{i \in R_r} x_{i,j} \geq 1 \quad \forall r \in [\rho] \qquad \text{(required courses are taken)}
$$

$$
\sum_{j=1}^{S} \sum_{i \in D} x_{i,j} \cdot u_i \ge 45
$$
\n
$$
\sum_{i=1}^{N} x_{i,j} \cdot y_{i,2-j} \le 1 \quad \forall t \in [T], \ \forall j \in [S]
$$
\n
$$
\sum_{k=1}^{\lceil S/2 \rceil} x_{i,2k-1} \le \sum_{t=1}^{T} y_{i,1,t} \quad \forall i \in [N]
$$
\n
$$
\sum_{k=1}^{\lfloor S/2 \rfloor} x_{i,2k} \le \sum_{t=1}^{T} y_{i,2,t} \quad \forall i \in [N]
$$
\n
$$
x_{i,j} \in \{0,1\} \quad \forall i \in [N], \ j \in [S]
$$

 $(45 \text{ units of depth}$ electives are taken)

 S (no classes are taken at the same time)

(no Spring-only course scheduled for Fall)

(no Fall-only course scheduled for Spring)

3 Implementing ILP with CMU-Specific Data

3.1 Data Sources

Our data for the project consisted of both FCE Data and Course Scheduling Data. The FCE data was obtained from the FCE database for the following semesters: Fall '21, Spring '22, Fall '22, Spring '23. The category "Hrs Per Week" was used as a measure for how hard and time consuming a class is. We averaged the FCE for each course across the four semesters in which data was collected from.

Our Scheduling Data was obtained from the CMU Schedule of Classes. We obtained the days and times that courses are offered for the Spring and Fall semesters (Fall '23, Spring '24 as that is what is currently viewable on the website). Furthermore, for each of our courses, we looked it up in the Schedule of Classes and obtained any prerequisite requirements. Here, we also obtained the actual units for each course. We formatted our data using Google Sheets, and transforming it into the necessary formats (such as for date/time of course offering) following the outline mentioned in our constraints (Section 2.3) for the ILP.

We selected classes that are required in the core curriculum as well as math, computer science, business, and statistics courses that can count towards the required depth electives for the major. We also included general education requirement courses that are outlined in the next section. There were a total of 96 courses that data was retrieved for.

3.2 Assumptions

In order for our ILP to output an accurate schedule in a reasonable amount of time, we had to make some initial assumptions. First, we assumed that the Schedule of Classes for Fall '23 and Spring '24 will be representative of the course offerings for all four years. This assumption is reasonable, as courses tend to be offered at the same time every year. Furthermore, we had to encode generic general education (Gen-Ed) requirement courses in order to meet the minimum number of units to graduate. To do so, we encoded 9 unit Gen-Ed requirements with various FCE units (between 5 and 9). We included nine of these courses at various times in both semesters. From our experience, typical Gen-Ed courses that students take vary between 5 and 9 FCE units and are offered various times and days. Thus, we believe these assumptions to accurately represent course offerings at CMU.

We assumed that the students come to CMU with AP credit for both Calculus I and Calculus II, as this is the case for most mathematics majors. There is a general outline for when courses should be taken, including first-year courses and the Engage requirements. We hard coded the classes in the the table below in order to follow the suggested timeline to complete the Engage requirements, as well as to complete first year requirements and necessary coursework (such as 21-128 and 38-101). It would be atypical for a student to take these courses outside of the suggested timeline, thus we fixed them for all outputs as follows:

Lastly, there are special topics statistics courses (36-46X) that vary in topics and offering times across semesters. For clarity and accuracy, we combined all of the special topics courses into one general special topics course, by averaging the FCEs for all of the offerings in the data. We selected at time in the spring and fall semesters that is accurate based on past data. Therefore, potential solutions will contain a general offering in which the student can choose from the semester-specific topic offerings.

4 Solving with Gurobi

For solving the ILP, we chose to use Gurobi Optimization Software. Though the software requires a purchased license to use, we had previously obtained an academic license from our 21-292 Operations Research I class. Thus, due to our familiarity with the software and availability, we chose to use it for the project. We coded our ILP using Python (code can be seen at the end of the paper) and the Gurobi software. We used the Python library Pandas in order get the data into Python and convert it into the necessary data structures for the model. When run locally on a computer from 2020, it took approximately 2 hours to produce a solution.

5 Solving with Heuristic Algorithm

In addition to the Gurobi optimization, we also created a heuristic algorithm in Python. Given course schedules with corresponding units and FCEs, course prerequisites, and graduation requirements, the algorithm outputs a feasible solution. The algorithm followed the same assumptions as the Gurobi model's, with one additional: each of the first six semesters has five classes total. Three of the classes must be technical and two are nontechnical.

After reading in the data files for each course, we created a dictionary with the class as the key and a list of the prerequisites needed in order to take said class as the value. Similarly, we created a second dictionary with each course, k , as a key, where the value was a list of classes whose prerequisites course k can be used to fulfill. Because there were a limited number of technical classes under consideration, a list of them was manually included. The next step included categorizing the classes which satisfied various graduation requirements, such as depth electives or taking a calculus course. Lastly, several boolean variables were initialized to indicate whether a specific requirement had been satisfied yet along with variables for the number of semesters to create a schedule for and the current number of depth and overall units.

We decided to implement a greedy algorithm in order to determine which classes to add to the current semester schedule. For each semester, among the classes that were not yet chosen, our algorithm picked the class that was the prerequisite for the most number of classes using our dictionary that we initialized above. For the class that was picked, the algorithm checked if all necessary prerequisites for the course were satisfied as well if there a section offered at the same time as a free slot in the schedule. The final condition was checking if the course was technical/nontechnical and if so, had the maximum three technical/two nontechnical condition had been satisfied yet. If these conditions were both satisfied, we added this course to our schedule and updated the graduation requirement indicator variables. We continued

this "picking" process until five classes were added to the current semester schedule.

It is important to note that a couple of courses were hard-coded due to assumptions on the model. First-year writing, Eureka, and 21-128: Concepts of Mathematics, were all added to the first semester. Similarly, the MCS Engage Requirements was added manually to each schedule, as was specified previously in Section 3.2.

The final output of this algorithm was a feasible course schedule.

6 Analyzing Results

In this section, we present a comparative analysis of our results with respect to various benchmarks, including:

6.1 Results from Gurobi

We first present the results obtained using Gurobi. These results serve as a foundational comparison for our findings.

Overall, the schedule is correct. For each course, all of the prerequisites have been satisfied in prior semesters and all of the graduation requirements have been satisfied. Moreover, the average number of FCEs for each semester, is within a similar range of 39 hours. Because the FCEs for each semester is within a similar range, this solution helps achieve the objective of minimizing the maximum FCE among the 8 semesters. We can see that the optimal schedule is extremely similar to our personal experiences so far with a healthy balance of technical and non-technical courses added to each semester.

One interesting thing to note in the schedule is that 38-304: Reading and Writing Science, is scheduled for Spring 1. Although this is a spring-only course, it is typically taken during Spring 3 rather than Spring 1. However, this is not a significant issue because it is not required, but recommended to be taken in Spring 3. Additionally, in Spring 4, 73-103: Principles of Macroeconomics and 73-230: Intermediate Microeconomics are both scheduled. This is not an issue either because 73-103 is not a prerequisite for 73-230, but students do not typically take these courses in the same semester.

6.2 Comparison to Heuristic Algorithm's Schedule

We want to compare the Gurobi optimal solution with the feasible one produced by the heuristic algorithm.

The feasible schedule using the greedy approach is as follows:

Compared to the optimal schedule, the feasible schedule had higher FCEs for all eight semesters, but were still relatively close to one another with the exception of Fall 3. The higher FCEs are likely due to the additional constraint of three technical and two nontechnical courses for each semester. Because of this assumption, more math classes were assigned to the schedule. Math classes typically have higher FCEs than the average nontechnical

course. Interestingly, the same number of nontechnical classes were assigned to both of the schedules, although not in the same semesters.

If planning a schedule that would allow you to have the most number of course options in following semesters, the heuristic algorithm may work better by choosing courses based on the number of courses it is a prerequisite for, especially in the earlier semesters. However, for a college experience with the work relatively evenly distributed among the eight semesters, the optimal solution is a possible way of achieving this goal while also ensuring that all of the graduation requirements are satisfied in time.

6.3 Comparison to Past Papers' Schedules

Next, we compare our derived schedules with those documented in prior papers. This comparative analysis highlights the deviations, improvements, or similarities between our schedules and those previously published.

There was a similar study from Fall 2019 where the students also scheduled courses for a mathematical sciences major. Their goal was to plan a schedule for a general math major that minimized the maximum number of FCEs and minimized the hours of free time during a student's week. Essentially, they wanted to a schedule that had the least amount of free time between classes as possible, so students can take classes back to back rather than having many small blocks of unproductive free time. Their results are as follows:

The maximum FCE units per semester using this schedule is 33.89 units

Total time between classes weekly for all semesters is 54 hours

Their schedule was mostly focused on scheduling only the math classes (21-xxx) over 8 semesters and when to take general electives. However, we chose to schedule all of the courses an Operations Research and Statistics major is required to take, besides the general electives. These included math, statistics, economics, accounting, computer science, and the MCS core courses such as the Engage requirements, the physical science requirement, the life science requirement, etc. The order of math classes were similar in the beginning, but began to diverge at semester 3/4. The previous students also allowed graduate courses in their schedule, which could be biased in terms of FCEs since grad students tend to spend less time on grad courses than undergrads do. They were also able to achieve a lower maximum FCE unit per semester of 33.89 compared to our maximum FCE unit of 38.8, although their solution doesn't meet the 360-unit graduation requirement.

There was another similar study conducted in Fall 2018. Their goal was to plan a schedule for a general math major, allowing studemts to choose from a list of depth electives that give the mathematics degree some "value" and take the easiest general education electives. They used a valuation system by using the overall scores from FCEs to evaluate the depth electives based on their value. The students used Integer Programming to determine which depth electives to take and a greedy algorithm to choose the general education electives that required the lowest FCEs. It should be noted that they assumed that all the courses a student wishes to take for a particular semester does not have any scheduling conflicts. Their results are as follows:

Their schedule was also mostly focused on scheduling only the math classes (21-xxx) over the four years and didn't include when to take general electives. They also planned the schedules by years, indicating which classes to take in a certain year, rather than a certain semester. As stated earlier, we were able to to schedule all of the courses an Operations Research and Statistics major is required to take, besides the general electives. There were also a few feasibility concerns with their outputted schedule, even when considering each "semester" as denoted in their table as a year. For example, 36-225 is a fall-only class and requires 21-120 as a pre-req, but both are under the same academic year. 21-292 is recommended to be taken in year 1, but requires 21-241 and 21-228 which are both scheduled in year 2. Moreover, 21-270 is a spring-only class and is a pre-req for 21-370, but both are scheduled in year 2. The order of math classes were similar in the beginning, but began to diverge at semester 3/4. This may be because their solution model used the most recent catalog (2018) and major requirements, but the course catalog may have been updated in the recent years.

While prior studies focused primarily on math courses, our schedule prioritized a wider array of general education requirements requirements, omitting bias potentially introduced by prioritizing specific courses (such as grad classes) or overlooking critical prerequisites and class time constraints. This holistic approach ensured a well-rounded curriculum for an Operations Research and Statistics major. By addressing these broader considerations, our schedule was not only feasible but also expanded the scope, enabling a more comprehensive optimized schedule for future students pursuing this major

6.4 Comparison to Published Schedule

Furthermore, we will compare our schedule against the CMU suggested schedule on the CMU website. The suggested schedule is as follows:

While sharing a similar pathway in math course scheduling, our analysis revealed a contrasting spread of FCEs over the suggested four-year period. Their recommended plan exhibited a wider range of FCEs, fluctuating between 36.5 to 50.4 hours per week. In contrast, our optimized schedule showcased a more consistent distribution, ranging from 38.5 to 38.8 hours per week. In other words, our schedule was able to minimize the variance in FCEs, ensuring a more balanced and manageable workload for students. Notably, our selection of depth electives aimed at minimizing the maximum FCE, instead of leaving these choices to individual student discretion (like the recommended schedule did). Additionally, our ordering of non-math courses, including economics and accounting, differed.

7 Conclusion

Our study in optimizing college courses has yielded interesting observations. We saw that our schedule managed to evenly distribute the workload across all semesters. This resulted in a balanced academic course load, ensuring students wouldn't face overwhelming spikes in coursework at any point during their college experience. The optimized schedule was also very similar to our own academic paths and mirrored the courses many of us have taken or planned to take. This correlation emphasizes the relevance and applicability of our findings to real-life academic journeys, further validating the importance of thoughtful course planning and optimization for a smoother college experience.

However, our study isn't without its limitations. The dynamic nature of class schedules and the variability in FCEs due to individual professors pose challenges in creating a universally applicable scheduling strategy.

Broadening this approach to encompass a variety of majors, minors, and concentrations beyond Operations Research and Statistics could offer valuable insights into diverse academic schedules. We could also consider optimizing the course planning strategy to minimize the overall FCE over the four-year period. Prioritizing classes that might increase workload but significantly contribute to future job prospects, such as those in machine learning, finance, or programming, could be potentially beneficial for students as well. Another modification could be to minimize the FCE during senior year, so students can have more time to apply to jobs and/or graduate schools.

Our study holds practical significance in addressing the challenge of managing academic loads in college effectively. Enhancing the scheduling process could help students in balancing their academics with other non-academic interests such as extracurriculars, sports, etc.

8 Appendix: Code for Gurobi Solver

The following four pages show the Python code used to encode our ILP and run it through the Gurobi optimization software:

```
from gurobipy import *
   #------------------Defining the model---------------------------
 # Initialization. The name is arbitrary
4
  model = Model('take3')
  #-----------------Importing Data-----------------------------
 import pandas as pd
9
import numpy as np
10
sched = pd.read_csv("PD_Schedule.csv")
12
units = pd.read_csv("PD_Units.csv")
13
PR = pd.read_csv("PD_PR.csv")
14
groups = pd.read_csv("PD_Groups.csv")
15
req = pd.read_csv("PD_Req.csv")
16
semCount = 8 # semCount is the number of semesters
18
courseCount = len(units) # courseCount is the number of courses
19
# Schedule Matrix
21
fall = sched[sched.Sem == "F"]
22
fall = fall.drop("Sem", axis = 1)
23
fall = fall.drop("Num", axis = 1)
24
spr = sched[sched.Sem == "S"]
26
spr = spr.drop("Sem", axis = 1)
27
spr = spr.drop("Num", axis = 1)
28
timeVec = list(np.stack((fall, spr), axis = 1)) 
30
#timeVec_i,s,g is if course i is offered in timeslot g in season s (fall or
31
   spring)
timeCount = len(timeVec[0][0])
33
34 | print(timeCount)
# FCE Vector 
36
fceVec = list(units["FCE"]) # f_i is the FCE of course i
37
# Units Vec 
39
unitVec = list(units["Units"]) # unitVec_i is the units of course i
40
# Group Vec 
42
groups = groups.drop("Num", axis = 1)
43
groupVec = groups.values.tolist()
44
groupCount = len(groupVec[0])
45
   "''''"47 groupVec layout (group1 contains course 1 and n)
                         Group 1     Group 2     ...     Group n
49 Course 1 [1 0 0 0]
50 Course 2 [0 0 0 1]
...
51
Course n   [1             0                   0]
52
   """
 1
 2
3
 5
6
7
8
11
17
20
25
29
32
35
38
41
46
48
53
54
```

```
# PR Vec 
56
 PR = PR.drop("Num", axis = 1)
57
 prereqVec = PR.values.tolist()
58
     prereqCount = len(prereqVec[0])
 # says the group numbers that are prereqs for each course
60
     "''''"''"prereqVec layout (course 1 has prereq of group 2)
                                Group 1     Group 2     ...     Group 3
     \begin{bmatrix} \text{Course 1} & \begin{bmatrix} 0 & & & 1 \end{bmatrix} & \begin{bmatrix}\begin{bmatrix} \text{Course 2} & \text{[0]} & \text{0} & \text{0} & \text{0} \end{bmatrix}...
 Course n   [0             1                   0]
67
     "''''"''"# Requirements Vector 
71
 req = req.drop("Num", axis = 1)
72
 depthVec = req["Depth Electives"]
73
 depthReqUnits = 45
74
 req = req.drop("Depth Electives", axis = 1)
75
 reqVec = req.values.tolist()
76
 reqCount = len(reqVec[0])
77
 # says the courses needed to fulfill requirments 
78
      "''"reqVec layout (to fulfill requirement 1 need course 1 or 2)
                      \text{Req } 1 \qquad \text{Req } 2 \qquad \ldotsCourse 1   [1         0               
82
 Course 2   [1         0                
83
     ...
 Course n   [0         1         
85
     """
               #-----------------Creating decision variables---------------------
     # Define binary variables x_i,j
 x=[([0]*semCount) for i in range(courseCount)];
95
 for i in range(courseCount):
96
              for j in range(semCount):
                        x[i][j] = model.addVar(vtype=GRB.BINARY, 
                                                                               name="x_({},{})".format(i+1, j+1))
a = model.addVar(vtype=GRB.CONTINUOUS, name = "semester_max")
101
     # Pushing created variables to the model
model.update()
104
# Each class is taken once
106
for i in range(courseCount):
107
              model.addConstr(sum(x[i][j] for j in range(semCount)) <= 1, 
                                                    name='Class {} taken once'.format(i+1))
 59
 61
 62
 63
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 80
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 84
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 87
 88
 89
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103
105
108
109
```

```
# Don't overload in any semester
111
for j in range(semCount):
112
             model.addConstr(sum(unitVec[i]*x[i][j] for i in range(courseCount)) <= 54, 
                                                name="Don't overload in sem {}".format(j+1))
# Loop through courses
116
# Loop through groups
117
# If group is prereq for that course
118
# loop through group, add constr
119
for i in range(courseCount):
120
             for g in range(groupCount):
               if (prereqVec[i][q] == 1):
                      # If course i requires group g as a prereq
                               for j in range(semCount):
                                       # Then if course i is taken in semester j, at least one class (c)
                                       # from group g must be taken in a semester (k) lower than j
                                       model.addConstr(sum(sum(groupVec[c][g] * x[c][k] for k in
     range(j)) 
                                                                                   for c in range(courseCount)) >= x[i][j], 
                                                                                   name="Group {} prereq for {} in sem
     {}".format(g+1,i+1,j+1))
# Each graduation requirement must be satisfied
131
for r in range(reqCount):
132
             model.addConstr(sum(sum(reqVec[i][r] * x[i][j] for i in range(courseCount)) 
                                                         for j in range(semCount)) >= 1, 
                                                         name="Satisfy requirement {}".format(r+1))
    # Need at least 45 units of depth electives
     model.addConstr(sum(sum(depthVec[i]*x[i][j]*unitVec[i] for i in
     range(courseCount)) 
                                                for j in range(semCount)) >= depthReqUnits,
                                                name = "Satisfy depth requirements")
# One class per time slot per semester
142
for g in range(timeCount):
143
             for j in range(semCount):
                      model.addConstr(sum(x[i][j]*timeVec[i][j%2][g] 
                                                                  for i in range(courseCount)) <= 1, 
                                                                  name="One class in slot {} for sem
     {}".format(g+1,j+1))
# Courses only taken in seasons they are offered in
149
for i in range(courseCount):
150
             model.addConstr(sum(x[i][2*j] for j in range(4)) <=
                                                sum(timeVec[i][0][g] for g in range(timeCount)), 
                                                name="Course {} not taken in fall if spring only".format(i+1))
             model.addConstr(sum(x[i][2*j+1] for j in range(4)) <=
                                                sum(timeVec[i][1][g] for g in range(timeCount)), 
                                                name="Course {} not taken in spring if fall only".format(i+1))
# 360 units in total needed to graduate
159
model addConstr(sum(sum(x[i][j] * unitVec[i] for j in range(semCount))
160113
114
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158
```

```
for g in range(timeCount):
        for j in range(semCount):
                    model.addConstr(sum(x[i][j]*timeVec[i][j%2][g] 
                                                               for i in range(courseCount)) <= 1, 
                                                               name="One class in slot {} for sem
    {}".format(g+1,j+1))
# Courses only taken in seasons they are offered in
149
for i in range(courseCount):
150
            model.addConstr(sum(x[i][2*j] for j in range(4)) <=
                                              sum(timeVec[i][0][g] for g in range(timeCount)), 
                                              name="Course {} not taken in fall if spring only".format(i+1))
            model.addConstr(sum(x[i][2*j+1] for j in range(4)) <=
                                              sum(timeVec[i][1][g] for g in range(timeCount)), 
                                              name="Course {} not taken in spring if fall only".format(i+1))
# 360 units in total needed to graduate
159
    model.addConstr(sum(sum(x[i][j] * unitVec[i] for j in range(semCount)) 
                                              for i in range(courseCount)) >= 360, 
                                              name = "Total courses needed to graduate")
# Hardcoding to make schedule more reliable
165
x[44][0] = 1; # concepts in sem 1
166
x[72][0] = 1 # eureka in sem 1
167
x[75][3] = 1; # engageInwards in sem 4
168
x[76][4] = 1; # engageOutwards in sem 5
169
x[77][6] = 1; # engageForwards in sem 7
170
x[73][6] = 1; # enagageService in sem 7
171
x[74][6] = 1; # engageArts in sem 7
172
x[71][2] = 1 # underGradColloq in sem 3
173
model.addConstr(sum(x[65][j]+x[67][j]+x[68][j]+x[69][j] for j in range(2,8)) <=0, 
174
                                      name="No first year writing after first year")
# To minimize the max FCE among semesters (but keep this linear), we create
178
# variable a which is at least the FCE for each semester and minimize that
179
model.setObjective(a, GRB.MINIMIZE)
180
for j in range(semCount):
181
            model.addConstr(a >= sum(fceVec[i] * x[i][j] for i in range(courseCount)), 
                                              name = "Max FCE >= sem {} FCE".format(j+1))
# Printing the model in a separate file for easier look
185
model.write('take3.lp')
186
    #-------------------- Solving the LP -------------------------------------
model.optimize()
191
#------------------ Outputting the solution ----------------------------
193
# Prints the non-zero variables and its values in a table format
195
model.printAttr('X')
196
\overline{1}144
145
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190
192
194
-
```
9 Appendix: Code for Heuristic Algorithm

```
2 '''
3 heuristic algorithm
4 input: course schedules
5 output : feasible schedule of classes
 6
7 ** assumption : 3 technicals , 2 non - technicals
8
9 1. read in course schedule
10 2. make dictionary based on prerequisites (class --> classes prereq for)
11 3. make dictionary (class \rightarrow classes needed as a prereq)
12 4. 8 times: loop and choose 3 technicals randomly --> add to list of
     visited classes
13 5. have a list for classes chosen in one semester
14
15 ''''
16
17 import pandas as pd
18 import numpy as np
19
20 # read in data
21 sched = pd . read_csv (" PD_Schedule .csv")
22 units = pd . read_csv (" PD_Units .csv")
23 prereq = pd.read_csv("PD_PR.csv")
24 groups = pd.read_csv("PD_Groups.csv")
25 req = pd . read_csv (" PD_Req .csv")
2627 needed_prereq = \{\}28 \text{prereq\_for} = \{\}29
30 # for each class , adds in prereqs needed
31 for i in prereq . index :
32 curr_class_i = prereq.loc[i]['Num']
33 classes_needed_for_i = [ col for col in prereq . columns if prereq . at [i ,
     col] == 1]
34
35 needed_prereq [ curr_class_i ] = classes_needed_for_i
36
37 prereq_classes = [('0R<sup>PR</sup>', 21292), ('DTF<sup>PR</sup>', 21370), ('MF<sup>PR'</sup>, 21270),
     ('DE_PR ', 21260) , ('Py_PR ', 15112) , ('D_PR ', 21228) , ('Con_PR ', 21128) ,
      ('Cal_PR ', 21268) , ('M_Pr ', 21241) , ('Al_PR ', 21373) , ('An_PR ', 21355)
      , ('P_PR ', 21325) , ('MI_PR ', 73102) , ('MA_PR ', 73103) , ('OM_PR ', 70371)
      , ('S_PR ', 36226) , ('MR_PR ', 36401) ]
38
39 # for each class , adds in what classes require it as a prereq
40 for c in prereq_classes :
```

```
41 curr_prereq = c [0]
42 filtered_classes = prereq [ prereq [ curr_prereq ] == 1]
43
44 curr_num = c [1]
45 prereq_for [ curr_num ] = filtered_classes ['Num ']. tolist ()
46
47 technicals = [21228 , 21236 , 21238 , 21241 , 21254 , 21256 , 21259 , 21260 ,
      21261 , 21266 , 21268 , 21269 , 21270 , 21292 , 21301 , 21325 , 21329 , 21341 ,
      21355 , 21356 , 21366 , 21369 , 21373 , 21374 , 21378 , 21380 , 21420 , 21441 ,
      21484 , 15110 , 15112 , 15122 , 15150 , 15210 , 21128 , 21237 , 21240 , 21242 ,
      21300 , 21370 , 21371 , 21393 , 21469 , 21360]
48
49 # classes that satisfy each of these requirements
50 depth_req = (\text{req}[\text{req}]) req [\text{length}]\text{ Electives} = (1) (\text{Num'})\text{ .} tolist ()
51 d_{\text{req}} = (req[req['D_R'] == 1]) ['Num'].tolist()52 \text{ m }req = (req[req['M_R'] == 1])['Num'].tolist()
53 cal_req = (req[req['Cal_R '] == 1])['Num'].tolist()
54 de_req = (req[req['DE_R'] == 1])['Num'].tolist()
55 p_{\texttt{req}} = (req \texttt{req} [ 'P_{R} ' ] == 1]) [ 'Num' ] . to list ()56 e_req = (req[req['E_R '] == 1])['Num'].tolist()
57 bs_req = (req[req['BS_R'] == 1])['Num'].tolist()
58 \text{ ps\_req} = (\text{req} [\text{req} [\text{'PS\_R'}] == 1]) [\text{'Num'}].\text{tolist()}59
60 # checks if requirements have been satisfied
61 depth_req_sat = False
62 d_req_sat = False
63 m_req_sat = False
64 cal_req_sat = False
65 de_req_sat = False
66 p_req_sat = False
67 e_req_sat = False
68 bs_req_sat = False
69 ps_req_sat = False
70 total_req_sat = False
71
72 all_available_classes = sched ['Num ']. tolist ()
73 taken_classes = []
74
75 # initializes number of units and semesters , minimum number of depth and
      total units
76 num_sem = 8
77 depth_req_units = 45
78 total_req_units = 360
79 max_num_tech = 3
80 max_num_non = 2
81
82 curr_depth_units = 0
83 curr_total_units = 0
```

```
85 final_schedule = {}
86
87 for s in range (1, num\_sem + 1):
88
89 curr_num_tech = 0
90 curr_num_non = 0
91
92 considered_classes = []
93
94 curr_schedule = [0] * 130
95
96 # initializes semesters with assumptions
97 if s = 1:
98 curr_sem_classes = [21128 , 38101 , 76101]
99 max_num_classes = 5
100 curr_num_tech = 1
101 curr_num_non = 2
102 elif s == 3:
103 curr_sem_classes = [21201]
104 max_num_classes = 6
105 elif s == 4:
106 curr_sem_classes = [38230]
107 max_num_classes = 6
108 elif s == 5:
109 curr_sem_classes = [38330]
110 max_num_classes = 6
111 elif s == 6:
112 curr_sem_classes = [38304]
113 max_num_classes = 6
114 elif s == 7:
115 curr_sem_classes = [38110 , 38220 , 38430]
116 max_num_classes = 8
117 else:
118 curr_sem_classes = []
119 max_num_classes = 5
120
121 curr_num_classes = len(curr_sem_classes)
122
123 while curr_num_classes < max_num_classes :
124 # greedily pick classes that is the prereq for most from technical
      + semester
125 # then check if prereqs needed are satisfied
126 # then check if any other classes from chosen are scheduled during
      that time
127 # then check if grad requirements are needed
128 prereq_sat = False
129 schedule_conflict = True
```

```
130 not_taken = False
131 not_considered = False
132 offered_in_sem = False
133
134 # gets set of classes that have not been taken yet or attempted to
      have been taken
135 filtered_dict = {key: value for key, value in prereq_for.items ()
     if (key not in considered_classes) and (key not in taken_classes)}
136 if not filtered_dict:
137 print ("The dictionary is empty.")
138 break
139
140 # finds class that the most number of classes require it
141 max_prereq_class = max (filtered_dict, key=lambda k: len(
     filtered\_dict[k]))
142
143 # checks if all prereqs for chosen class have been taken
144 if all (c in taken_classes for c in needed_prereq [str (
     max_prereq_class ) ]) :
145 prereq_sat = True
146
147 # finds what semester it is offered
148 index = sched.index [sched ['Num'] == str(max_prereq_class)].tolist
     () [0]
149 print (index)
150 print (sched)
151 print (sched ['Sem'])
152 curr_offered_sem = sched ['Sem '][ index ]
153
154 # if the current semester matches the semester when the course is
     offered
155 if ((s \; % \; 2) == 0) and curr_offered_sem == 'S') or (((s \; % \; 2) == 1)and curr_offered_sem == 'F'):
156 offered_in_sem = True
157
158 #check scheduling conflict
159 curr_offered_sem = sched ['Sem'] [index]
160
161 max_prereq_class_sched = sched . loc [ index , sched . columns . difference
     ([ 'Num', 'Sen'] )162
163 # check if there is an avaialble time in current schedule for when
      the class is offered
164 result = [(a + b) % 2 for a, b in zip(curr_scheduled)]max_prereq_class_sched ) ]
165
166 chosen_time = []
167 chosen_yet = False
```

```
169 for i in range (len (curr_schedule) - 1):
170 if not chosen_yet:
171 if result [i] == 1 and result [i + 1] == 1:
172 schedule_conflict = False
173 chosen_yet = True
174 chosen_time . append (i)
175 chosen_time . append (i + 1)176 if i != (len(curr_schedule) - 2):
177 if result [1 + 2] == 1:
178 chosen_time . append ( i + 2) chosen_time . append ( i + 2)
179
180 # checks if class has been considered or taken yet
181 if max_prereq_class not in considered_classes:
182 not_considered = True
183
184 if max_prereq_class in all_available_classes :
185 not_taken = True
186
187 # increments either nontechnical or technical count by 1
188 if max_prereq_class in technicals:
189 curr_num_tech += 1
190 else:
191 curr_num_non += 1
192
193 # checks if all constraints satisfied
194 if prereq_sat and (not schedule_conflict) and not_taken and (
    curr_num_tech <= max_num_tech ) and ( curr_num_non <= max_num_non ) and
    not_considered and offered_in_sem :
195
196 print (f"adding {max_prereq_class}")
197
198 curr_sem_classes.append (max_prereq_class)
199 all_available_classes . remove (max_prereq_class)
200 taken_classes . append ( max_prereq_class )
201 prereq_for . pop ( max_prereq_class )
202
203 index = units.index [units ] = max\_prereq_class].tolist ()
     [0]204 curr_total_units += units [index] ['Units']
205
206 for t in range (len (chosen_time)):
207 curr_schedule [chosen_time [t]] == 1
208
209 considered_classes . append ( max_prereq_class )
210
211 # count as grad req
212 if curr_total_units == total_req_units :
```

```
213 total_req_sat = True
214
215 if max_prereq_class in depth_req :
216 index = units.index [units] \times \frac{1}{216} = max_p = max_p = d.[0]
217 curr_depth_units += units [index] ['Units']
218
219 if curr_depth_units == depth_req_units :
220 depth_req_sat = True
221
222 if max_prereq_class in d_req :
223 d_req_sat = True
224 if max_prereq_class in m_req :
225 m_req_sat = True
226 if max_prereq_class in cal_req :
227 cal_req_sat = True
228 if max_prereq_class in de_req :
229 de_req_sat = True
230 if max_prereq_class in p_req :
p_{231} p_req_sat = True
232 if max_prereq_class in e_req :
233 e_req_sat = True
234 if max_prereq_class in bs_req :
235 bs_req_sat = True
236 if max_prereq_class in ps_req :
237 ps_req_sat = True
238
239 final_schedule [s] = (curr_sem_classes, curr_schedule)
240 print (final_schedule)
```