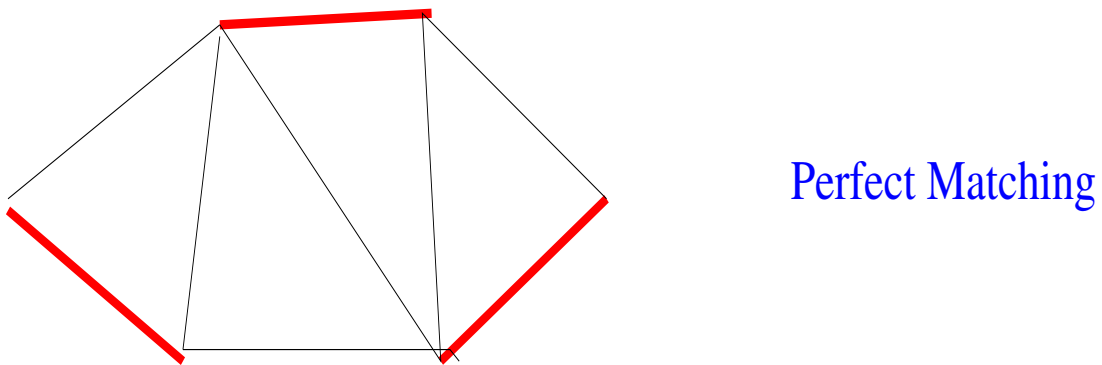
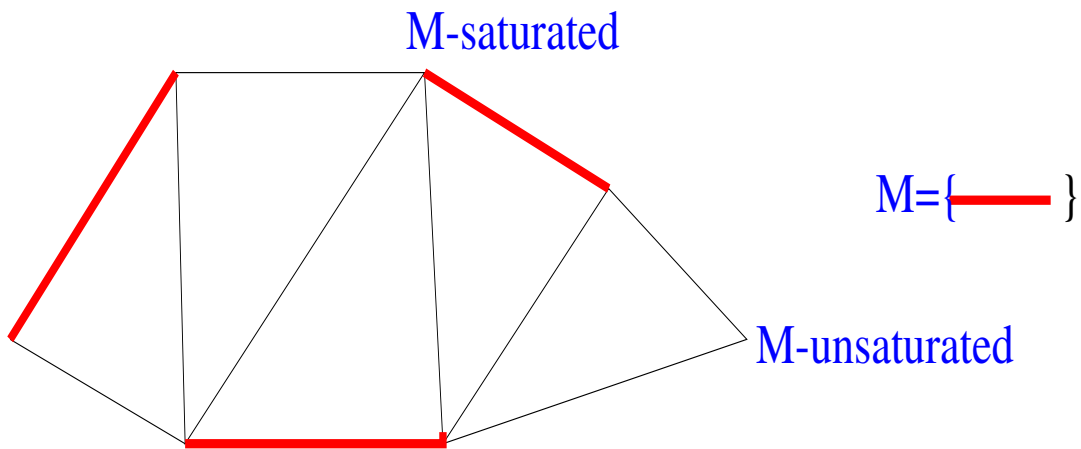
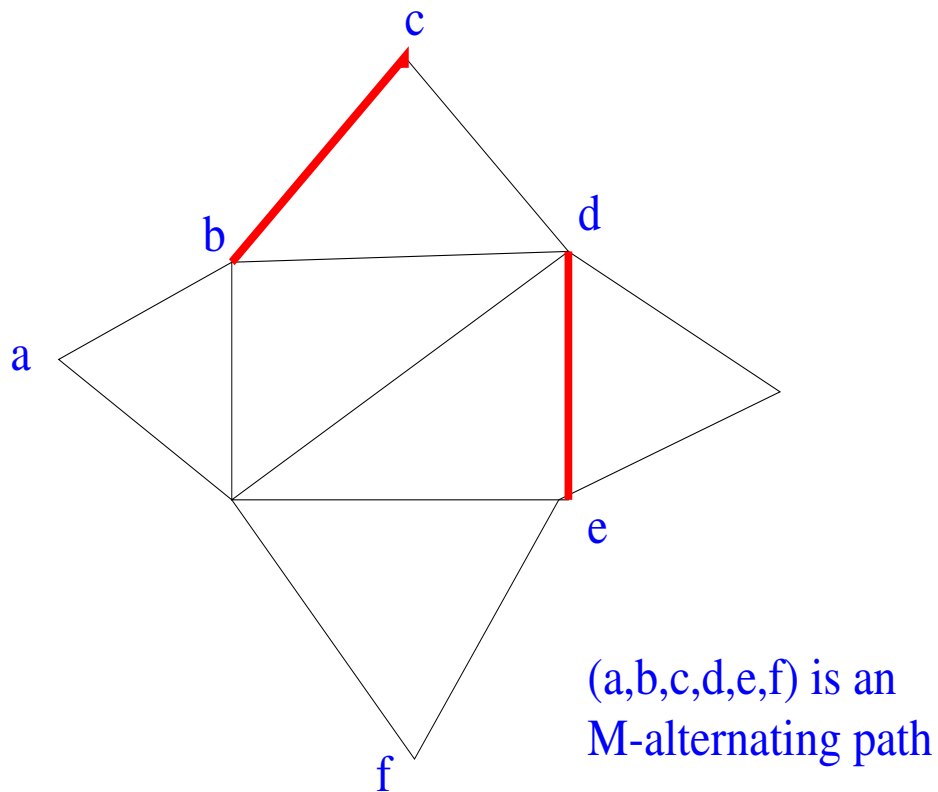


Matchings

A *matching* M of a graph $G = (V, E)$ is a set of edges, no two of which are incident to a common vertex.



M-alternating path

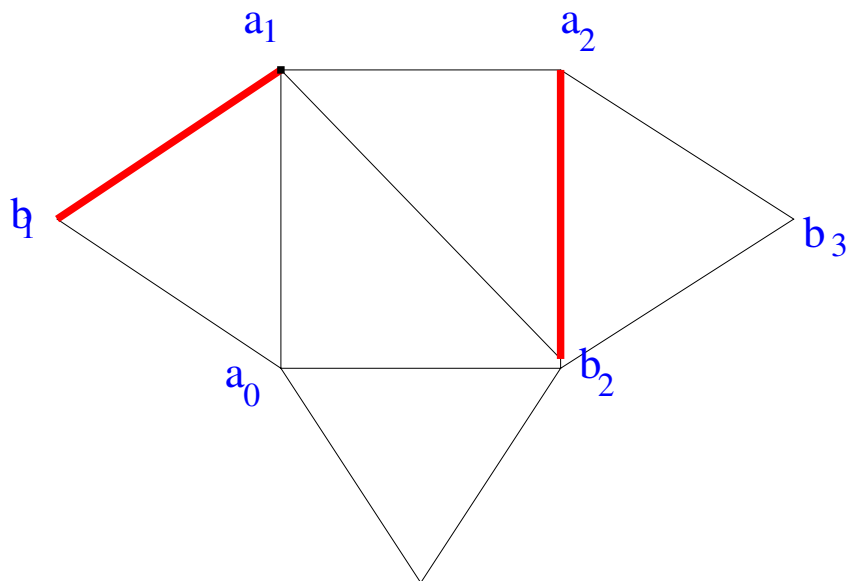


An M -alternating path joining 2 M -unsaturated vertices is called an M -augmenting path.

M is a *maximum* matching of G if no matching M' has more edges.

Theorem 1 M is a maximum matching iff M admits no M -augmenting paths.

Proof Suppose M has an augmenting path $P = (a_0, b_1, a_1, \dots, a_k, b_{k+1})$ where $e_i = (a_{i-1}, b_i) \notin M$, $1 \leq i \leq k+1$ and $f_i = (b_i, a_i) \in M$, $1 \leq i \leq k$.



$$M' = M - \{f_1, f_2, \dots, f_k\} + \{e_1, e_2, \dots, e_{k+1}\}.$$

- $|M'| = |M| + 1$.
- M' is a matching

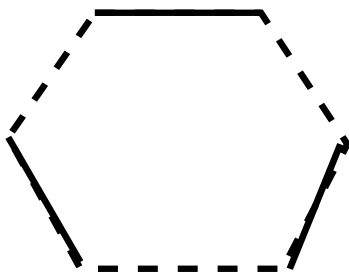
For $x \in V$ let $d_M(x)$ denote the degree of x in matching M , So $d_M(x)$ is 0 or 1.

$$d_{M'}(x) = \begin{cases} d_M(x) & x \notin \{a_0, b_1, \dots, b_{k+1}\} \\ d_M(x) & x \in \{b_1, \dots, a_k\} \\ d_M(x) + 1 & x \in \{a_0, b_{k+1}\} \end{cases}$$

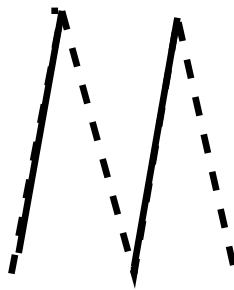
So if M has an augmenting path it is not maximum.

Suppose M is not a maximum matching and $|M'| > |M|$. Consider $H = G[M \nabla M']$ where $M \nabla M' = (M \setminus M') \cup (M' \setminus M)$ is the set of edges in *exactly* one of M, M' .

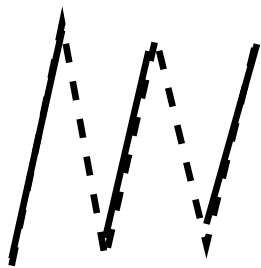
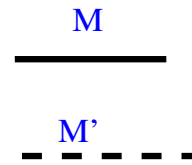
Maximum degree of H is 2 – ≤ 1 edge from M or M' . So H is a collection of vertex disjoint alternating paths and cycles.



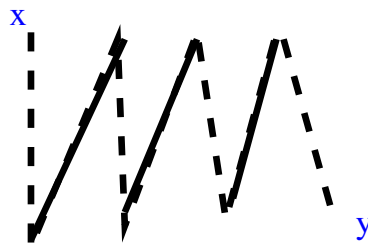
(a)



(b)



(c)



(d)

x, y M -unsaturated

$|M'| > |M|$ implies that there is at least one path of type (d).

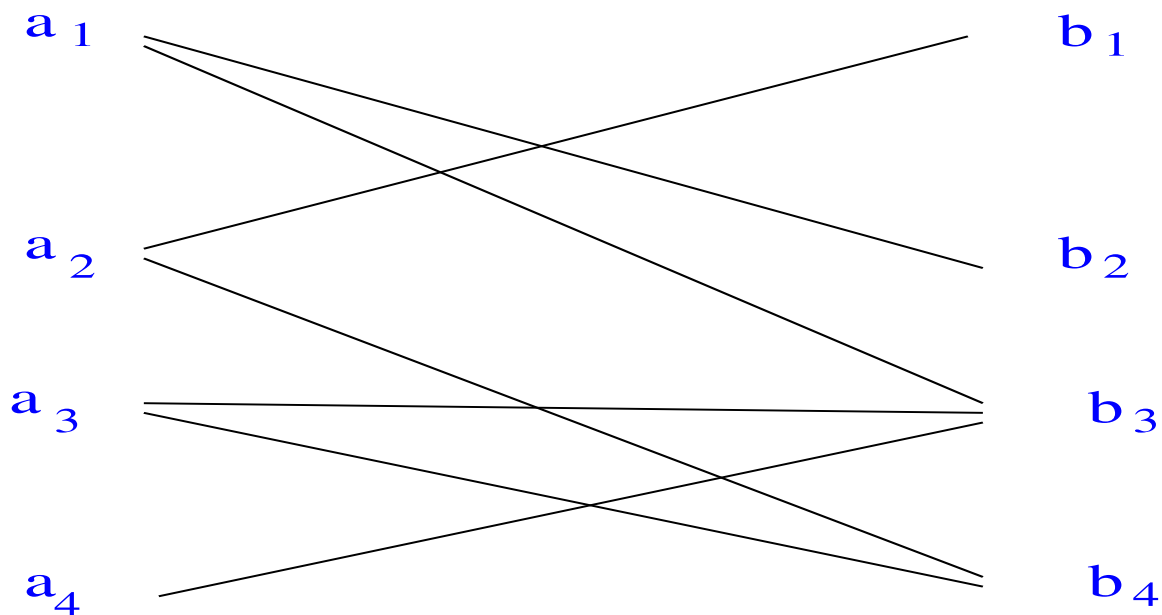
Such a path is M -augmenting

□

Bipartite Graphs

Let $G = (A \cup B, E)$ be a bipartite graph with bipartition A, B .

For $S \subseteq A$ let $N(S) = \{b \in B : \exists a \in S, (a, b) \in E\}$.



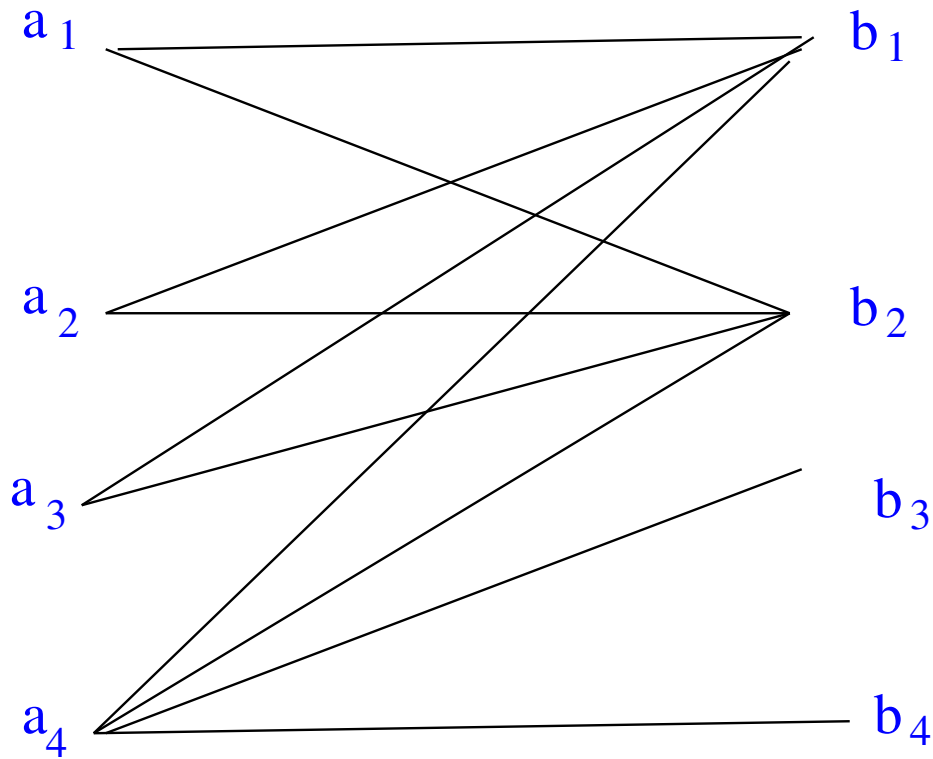
$$N(\{a_2, a_3\}) = \{b_1, b_3, b_4\}$$

Clearly, $|M| \leq |A|, |B|$ for any matching M of G .

Hall's Theorem

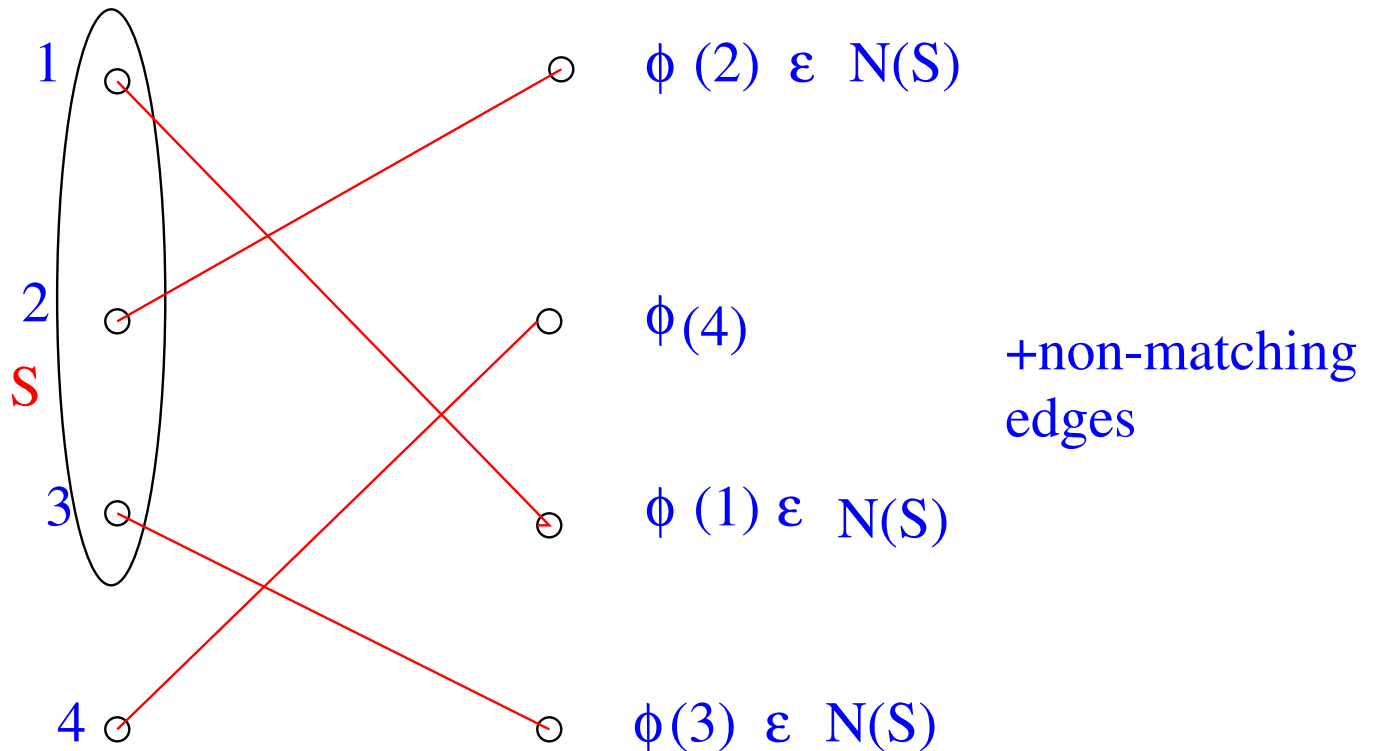
Theorem 2 G contains a matching of size $|A|$ iff

$$|N(S)| \geq |S| \quad \forall S \subseteq A. \quad (1)$$



$N(\{a_1, a_2, a_3\}) = \{b_1, b_2\}$ and so at most 2 of a_1, a_2, a_3 can be saturated by a matching.

Only if: Suppose $M = \{(a, \phi(a)) : a \in A\}$ saturates A .



$$|N(S)| \geq |\{\phi(s) : s \in S\}| \\ = |S|$$

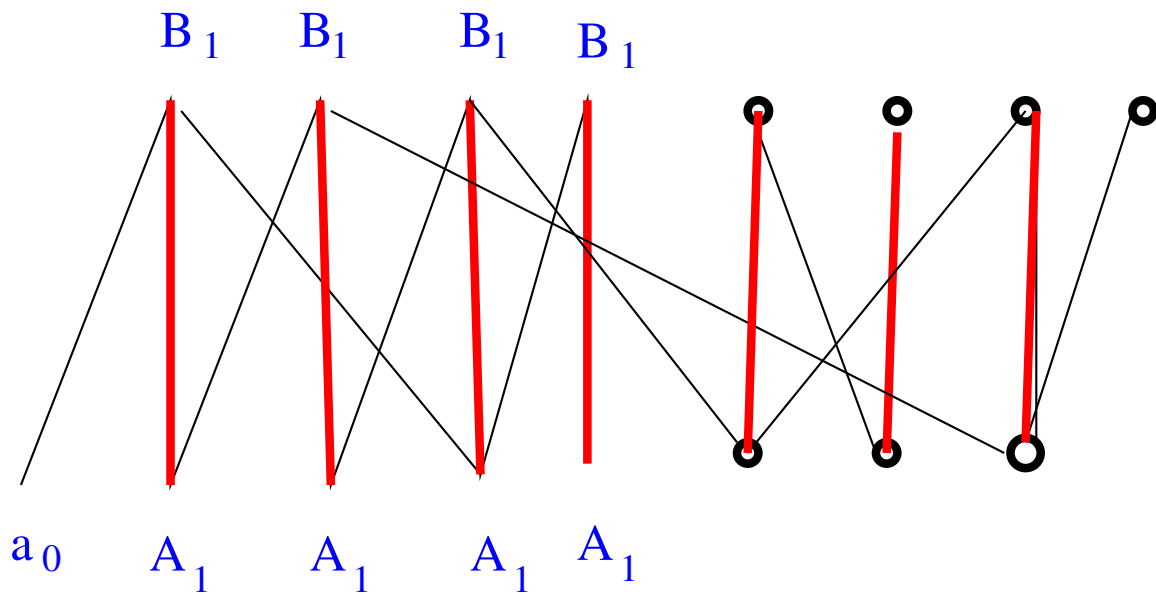
and so (1) holds.

If: Let $M = \{(a, \phi(a)) : a \in A'\}$ ($A' \subseteq A$) is a maximum matching. Suppose $a_0 \in A$ is M -unsaturated. We show that (1) fails.

Let

$A_1 = \{a \in A : \text{such that } a \text{ is reachable from } a_0 \text{ by an } M\text{-alternating path.}\}$

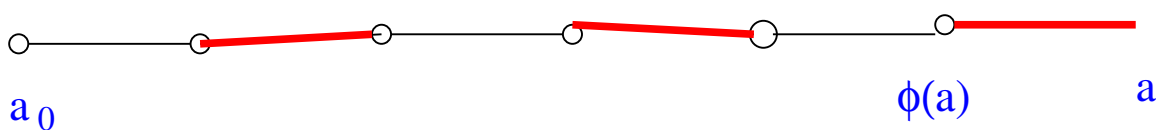
$B_1 = \{b \in B : \text{such that } b \text{ is reachable from } a_0 \text{ by an } M\text{-alternating path.}\}$



No $A_1 - B \setminus B_1$
edges

- B_1 is M -saturated else there exists an M -augmenting path.

- If $a \in A_1 \setminus \{a_0\}$ then $\phi(a) \in B_1$.

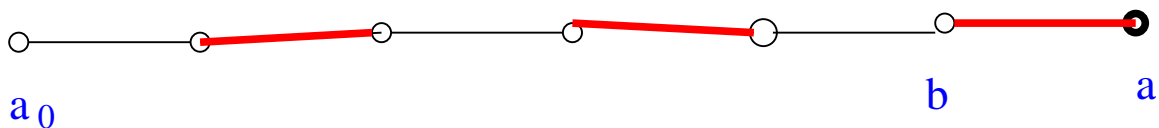


- If $b \in B_1$ then $\phi^{-1}(b) \in A_1 \setminus \{a_0\}$.

So

$$|B_1| = |A_1| - 1.$$

- $N(A_1) \subseteq B_1$



So

$$|N(A_1)| = |A_1| - 1$$

and (1) fails to hold.

Marriage Theorem

Theorem 3 Suppose $G = (A \cup B, E)$ is k -regular. ($k \geq 1$) i.e. $d_G(v) = k$ for all $v \in A \cup B$. Then G has a perfect matching.

Proof

$$k|A| = |E| = k|B|$$

and so $|A| = |B|$.

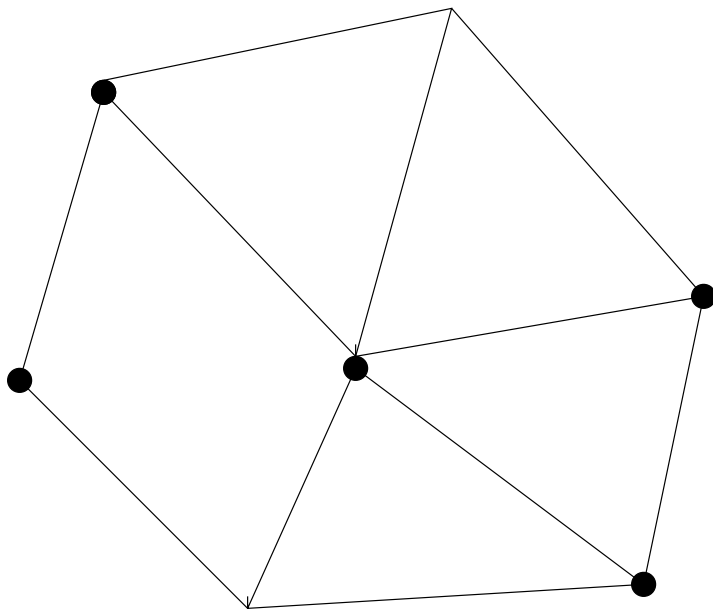
Suppose $S \subseteq A$. Let m be the number of edges incident with S . Then

$$k|S| = m \leq k|N(S)|.$$

So (1) holds and there is a matching of size $|A|$ i.e. a perfect matching.

Edge Covers

A set of vertices $X \subseteq V$ is a *covering* of $G = (V, E)$ if every edge of E contains at least one endpoint in X .



$\{\bullet\}$ is a covering

Lemma 1 If X is a covering and M is a matching then $|X| \geq |M|$.

Proof Let $M = \{(a_1, b_i) : 1 \leq i \leq k\}$. Then $|X| \geq |M|$ since $a_i \in X$ or $b_i \in X$ for $1 \leq i \leq k$ and a_1, \dots, b_k are distinct. \square

Konig's Theorem

Let $\mu(G)$ be the maximum size of a matching.

Let $\beta(G)$ be the minimum size of a covering.

Then

$$\mu(G) \leq \beta(G).$$

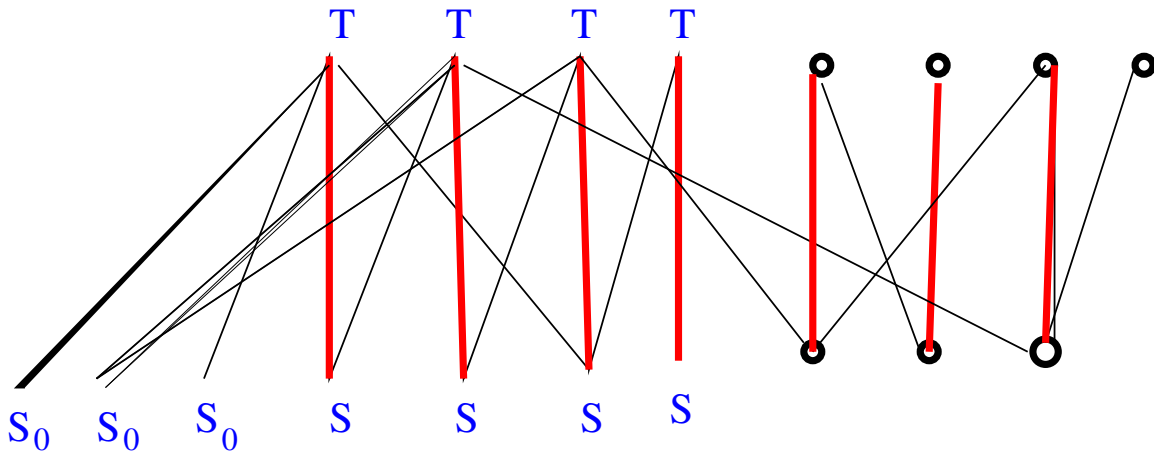
Theorem 4 *If G is bipartite then $\mu(G) = \beta(G)$.*

Proof Let M be a maximum matching.

Let S_0 be the M -unsaturated vertices of A .

Let $S \supseteq S_0$ be the A -vertices which are reachable from S_0 by M -alternating paths.

Let T be the M -neighbours of $S \setminus S_0$.



Let $X = (A \setminus S) \cup T$.

- $|X| = |M|$.

$|T| = |S \setminus S_0|$. The remaining edges of M cover $A \setminus S$ exactly once.

- X is a cover.

There are no edges (x, y) where $x \in S$ and $y \in B \setminus T$. Otherwise, since y is M -saturated (no M -augmenting paths) the M -neighbour of y would have to be in S , contradicting $y \notin T$. \square