

## Overview of the transportation algorithm

10 June

1. Form an initial basic feasible solution  $\{x_{ij}\}$  to the primal LP.

2. Use complementary slackness to find a corresponding solution  $\{v_i, w_j\}$  to the dual LP.

(By complementary slackness, when  $x_{ij}$  is basic, we should have  $v_i + w_j = c_{ij}$ .)

3. Compute test values  $c_{ij} - v_i - w_j$ .

(Test values for  $i, j$  corresponding to basic  $x_{ij}$  will always be zero, because of step 2. So we just need test values for nonbasic  $x_{ij}$ .)

4. If all test values are nonnegative, STOP: current solution  $\{x_{ij}\}$  is optimal.

(By second theorem.)

5. Otherwise, choose a negative test value (as a rule of thumb, the most negative one) and increase the corresponding  $x_{ij}$  as much as possible to decrease the cost. This is called a pivot. (Example shortly.) Then go back to step 2.

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## The transportation tableau

Rows correspond to origins.  
Columns correspond to destinations.

Example (same as previously):

Supply		Demand		Per-unit costs					
				TO					
A	10	W	8		W	X	Y	Z	
B	15	X	12	FROM	A	\$12	9	11	7
C	10	Y	7		B	15	8	10	9
		Z	8		C	7	4	6	11

Blank transportation tableau with this information:

	W:8	X:12	Y:7	Z:8
A:10	12	9	11	7
B:15	15	8	10	9
C:10	7	4	6	11

Squares in this grid correspond to variables in the LP, so we will have basic and nonbasic squares corresponding to basic and nonbasic variables.

First step: Get an initial basic feasible solution.

Method:

- Satisfy each demand one by one, from left to right.
- For each demand, take units from each supply, top to bottom, until the demand is met; but once a supply has been exhausted, it cannot be used to satisfy later demands.
- If necessary, add one or more basic squares with a value of zero. (More about this later.)

For this example, our initial basic feasible solution is:

	W: 8	X: 12	Y: 7	Z: 8
A: 10	8	2		
B: 15		10	5	
C: 10			2	8

Objective value  
(total cost):

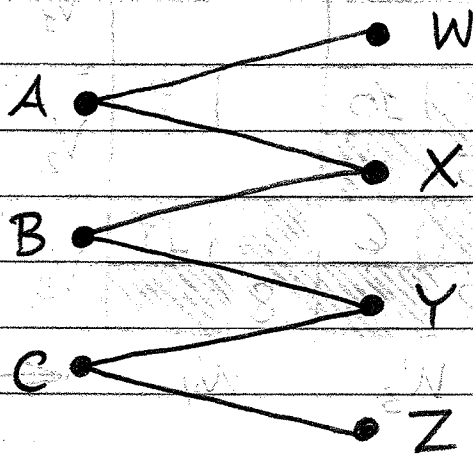
$$8(12) + 2(9) + 10(8) + 5(10) + 2(6) + 8(11) = \underline{\underline{\$344}}$$

Basic squares are shaded. Number in lower left corner is value of corresponding basic variable  $X_{ij}$ , i.e., number of units to be sent from origin  $i$  to destination  $j$ .

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## Transportation tableau — (2)

Note: Suppose we make a bipartite graph with vertices representing origins on the left and vertices representing destinations on the right, with an edge joining origin  $i$  to destination  $j$  if and only if the variable  $x_{ij}$  is basic:



Observe that this graph is a tree: it is connected (there is a path between any two vertices) and it has no cycles. (This particular graph happens to be a path graph, but that will not be true in general.)

This is an important fact. Basic solutions for the transportation problem will always correspond to trees.

★ If necessary, include one or more basic squares with value zero in order to get a tree.

Now that we have an initial basic feasible solution, we use complementary slackness to find a corresponding solution  $\{v_i, w_j\}$  to the dual LP. Complementary slackness implies that when  $x_{ij}$  is basic, we should have  $v_i + w_j = c_{ij}$ .

	W	X	Y	Z	
A	12	9	11	7	$v_1$
B	15	8	10	9	$v_2$
C	7	4	6	11	$v_3$
	$w_1$	$w_2$	$w_3$	$w_4$	

$v_i$ 's correspond to origins  
 $w_j$ 's correspond to destinations

From complementary slackness:

$$\begin{aligned}
 (x_{AW}) \quad v_1 + w_1 &= 12 \\
 (x_{AX}) \quad v_1 + w_2 &= 9 \\
 (x_{BX}) \quad v_2 + w_2 &= 8 \\
 (x_{BY}) \quad v_2 + w_3 &= 10 \\
 (x_{CY}) \quad v_3 + w_3 &= 6 \\
 (x_{CZ}) \quad v_3 + w_4 &= 11
 \end{aligned}$$

We have six equations in seven unknowns, so we have one free variable — we can set any one of these variables to any value we like, and then solve for the remaining six.

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### Transportation tableau — (3)

For simplicity, let's set  $v_1 = 0$ . Solving for the remaining six dual variables, we get

	W	X	Y	Z	
A	12	9	11	7	$v_1 = 0$
B	15	8	10	9	$v_2 = -1$
C	7	4	6	11	$v_3 = -5$

$w_1 = 12$     $w_2 = 9$     $w_3 = 11$     $w_4 = 16$

Now, for the nonbasic squares, we calculate test values,  $c_{ij} - v_i - w_j$ , and write them in the lower right corner:

	W	X	Y	Z	
A	12	9	11	7	$v_1 = 0$
B	15	8	10	9	$v_2 = -1$
C	7	4	6	11	$v_3 = -5$

$w_1 = 12$     $w_2 = 9$     $w_3 = 11$     $w_4 = 16$

We have negative test values, so this solution is not optimal. Choose the most negative ( $-9$  in AZ) and pivot there.

Pivoting. The test value  $-9$  in the AZ square indicates that if we make  $x_{AZ}$  positive (by bringing it into the basis) then we will decrease our total cost. So we want to increase the value of  $x_{AZ}$  from its current value of zero to some new value  $t$ .

But unless we adjust the values of the other  $x_{ij}$ 's, we will mess up our row and column sums. The total of the  $x_{ij}$ 's in the Z column must be 8 (to satisfy Z's demand), so if we increase  $x_{AZ}$  to  $t$  then we need to decrease  $x_{CZ}$  to the value  $8-t$ .

Now that adjustment messed up the total for the C row, so we need to increase the value of  $x_{CX}$  to  $2+t$ .

Following this chain of reasoning:

- (Y column) Decrease  $x_{BY}$  to  $5-t$ .
- (B row) Increase  $x_{BX}$  to  $10+t$ .
- (X column) Decrease  $x_{AX}$  to  $2-t$ .

This last adjustment also fixed the A row sum that we messed up by increasing  $x_{AZ}$  at the beginning. So now we're all good.



## Transportation pivoting - (2)

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Summary of these adjustments:

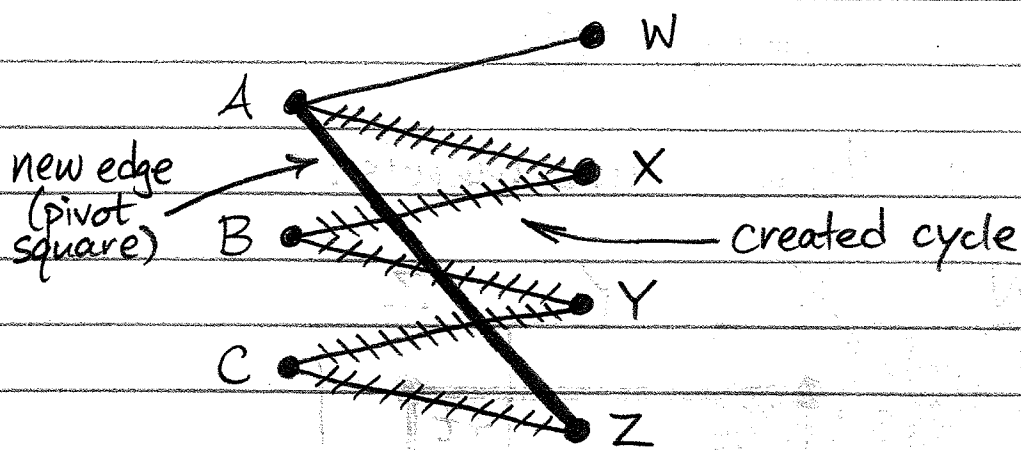
	W	X	Y	Z
A		$2-t$		$+t$
B		$10+t$	$5-t$	
C			$2+t$	$8-t$

Note that these adjustments follow a circuit of "alternating-direction rook's moves" (alternating vertical and horizontal motions) that, except for the pivot square AZ, turns  $90^\circ$  only on basic squares.

The adjustments alternate  $+t, -t, +t, -t, \dots$  at the "corners" of this circuit, starting with  $+t$  at the pivot square.

There will always be exactly one such circuit in the tableau for any given pivot square. This is because the basic squares correspond to the edges of a tree, and adding one edge (i.e., the pivot square) to a tree creates a unique cycle:





Next question: What should  $t$  be?

Every unit increase to the value of  $x_{AZ}$  will decrease our total cost by \$9 (that's what the test value means), so we want to make  $t$  as large as possible.

But if  $t$  is too large, the values of  $x_{AX}$ ,  $x_{BY}$ , and  $x_{CZ}$ , which are  $2-t$ ,  $5-t$ , and  $8-t$ , respectively, will become negative, which violates their domains. The first to become negative will be  $x_{AX} = 2-t$ , so the greatest we can make  $t$  is 2.

(Note that  $x_{BX} = 10+t$  and  $x_{CY} = 2+t$  do not place restrictions on  $t$ , because they will not become negative.)

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# Transportation pivoting - (3)

So, taking  $t=2$ , our new basic feasible solution is

↖ AX: fell out of basis ( $x_{AX}$  became zero)

	W	X	Y	Z		
A	8	12	9	11	7	
B		15	12	8	10	9
C		7	4	4	6	11

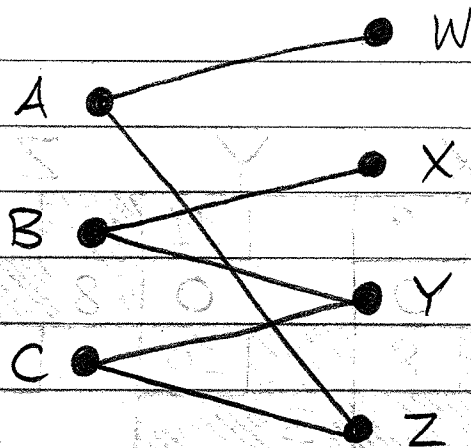
$x_{AW}$  is still 8 because AW was not part of the pivot circuit.

$AZ$ : new basic square

Objective value:

$$8(12) + 2(7) + 12(8) + 3(10) + 4(6) + 6(11) = \underline{\underline{\$326}}$$

(improved)



Still a tree.  
(connected, acyclic)

Compute dual variables and test values:

	W	X	Y	Z			
A	8	12	9	11	7	$V_1 = 0$ ← (arbitrariness)	
B		15	12	8	10	9	$V_2 = 8$
C		7	4	4	6	11	$V_3 = 4$

$W_1 = 12$     $W_2 = 0$     $W_3 = 2$     $W_4 = 7$

Still have negative test values, so still not optimal. Pivot on CW. Pivot circuit goes CW-AW-AZ-CZ-CW:

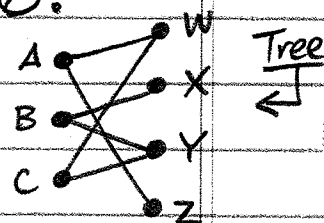
	W	X	Y	Z
A	8-t			2+t
B				
C	t			6-t

So largest possible value of  $t$  is 6.

Next tableau:

	W	X	Y	Z	
A	2			8	$V_1 = 0$
B	4			3	$V_2 = -1$
C	6			9	$V_3 = -5$

$$W_1 = 12 \quad W_2 = 9 \quad W_3 = 11 \quad W_4 = 7$$



All test values are nonnegative, so this solution is optimal:

2 units  $A \rightarrow W$ , 8 units  $A \rightarrow Z$ , 12 units  $B \rightarrow X$ ,  
3 units  $B \rightarrow Y$ , 6 units  $C \rightarrow W$ , 4 units  $C \rightarrow Y$ .

$$\text{Optimal objective value (cost)} = 2(12) + 8(7) + 12(8) + 3(10) + 6(7) + 4(6) = \underline{\underline{\$272.}}$$