

28 May

Combinatorial optimization overview

- Combinatorial optimization is the study of problems that involve a search for the "best" option among a (usually finite) set of choices.

- Example: Factory placement.

A company has ten possible sites at which it can build factories.

A factory can produce 100 units/month of output. Each site has a different building cost:

Site:	1	2	3	4	...
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Building cost:	\$1.2M	\$1.8M	\$0.8M	\$2.1M	...
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The company has a list of customers with monthly demands:

Customer:	A	B	C	D	E	...
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Monthly demand:	12	16	23	7	10	...
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Shipping costs per unit:

		FROM			...
		1	2	3	...

	A	\$4	\$6	\$13	...
TO	B	10	8	9	...

Where should factories be built to satisfy customer demand over 5 years at min cost?

	A	7	5	3	...
	B	7	5	3	...
	C	7	5	3	...

- Example: Bin packing.
 - List of items of various sizes:

Item:	1	2	3	4	...
Size:	3	7	12	9	...
 - Supply of bins, each having the same capacity (say, capacity 50).
 - How to place all items into bins, using the minimum number of bins, without exceeding bin capacity?
 - "Steel mill slab problem":
Each item has a color, and the number of different colors in each bin can be no more than 2.

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Combinatorial optimization overview

- Example: Swim meets. Two teams.
- Each team can enter at most three swimmers in any one individual (non-relay) event.
- Each team can provide at most three "entries" in each relay event, where an entry is a group of four swimmers.
- Each swimmer can enter at most four events.
- Each swimmer can enter at most two individual events.
- In a relay event, a team cannot be awarded points for more than two finishing places.
- Coach has information about event times of swimmers, and limited information about opposing team.
- How to assign swimmers to events to maximize probability of winning?

[Nowak, Epelman, Pollock]

More examples:

- Project scheduling (CPM)
- Minimum spanning tree
- Traveling salesman

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Linear programming — intro example.

- Farmer Brown is planning to raise potatoes and wheat.
- 100 acres of land available.
- 160 days of labor available.
- One acre of wheat will require 24 days of labor.
- One acre of potatoes will require 8 days of labor.
- \$1100 available for start-up costs of planting and cultivating.
- \$10 per acre to plant/cultivate potatoes.
- \$20 per acre to plant/cultivate wheat.
- Expected revenue:
\$40/acre for potatoes,
\$120/acre for wheat.
- How many acres of each crop should be planted to maximize total revenue?

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LP formulation.

1. Identify variables and their domains.

- Variables usually represent decisions to be made — quantities under your control.
- Imagine a form letter describing the best course of action to take, the best solution. Write a sentence that could be used to describe any solution, using blanks in place of numbers. Those blanks are the variables.
- Sometimes auxiliary variables can be helpful — this usually becomes evident when trying to formulate constraints or objective function.
- Each variable must have a corresponding domain. In an LP, there are three possible domains:

$$\underbrace{x \geq 0,}_{}$$

$$x \leq 0,$$

x unrestricted.

By far the most common.

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LP formulation.

2. Identify the objective.

- What quantity are you aiming to maximize or minimize?

3. Write the objective function in terms of the defined variables.

- Remember that the objective function must be linear.

4. Identify the constraints in the problem and express each one in terms of the defined variables.

- Remember that every constraint must be a linear inequality (or equality). Inequality cannot be strict.

— Linear means:

- No raising variables to powers,
- No multiplication of variables together,
- No dividing by variables,
- No functions like $\sin()$, $\sqrt{\cdot}$, etc. applied to variables.

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LP formulation for Farmer Brown example.

Variables:

p : number of acres of potatoes

w : number of acres of wheat

Domains: $p \geq 0, w \geq 0$

Objective: Maximize revenue.

Objective function: $40p + 120w$

Constraints:

Land (resource) : $p + w \leq 100$

Labor (resource) : $p + 4w \leq 160$

Capital (resource) : $10p + 20w \leq 1100$

Written : maximize $40p + 120w$

Subject to $p + w \leq 100$ [land]

$p + 4w \leq 160$ [labor]

$10p + 20w \leq 1100$ [capital]

$p \geq 0, w \geq 0$.

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"Anatomy" of an LP.

Objective →

$$\max \quad 40P + 120W$$

Objective
function

Constraints →

s.t.

$$\begin{aligned} P + W &\leq 100 \\ P + 4W &\leq 160 \\ 10P + 20W &\leq 1100 \end{aligned}$$

Variable domains →

$$P \geq 0, \quad W \geq 0$$

Terminology

- Solution: An assignment of values to variables.
- Feasible solution: A solution that satisfies all constraints (and domains).
- Feasible region (a.k.a. feasible set): The set of all feasible solutions.
- Objective value: The value of the objective function corresponding to a given solution.
- Optimal (feasible) solution: A feasible solution whose objective value is at least as "good" as that of any other feasible solution.
- Optimal objective value: The objective value of an optimal feasible solution.

More terminology.

- Feasible LP: An LP with at least one feasible solution.
- Infeasible LP: One that is not feasible.
- Unbounded LP: A feasible LP with no optimal feasible solution.

Solving an LP means attempting to find an optimal feasible solution (not just the optimal objective value!)

Possible outcomes:

- LP is infeasible.
- LP is feasible.
 - LP is unbounded.
 - LP has an optimal feasible solution.
 - LP has a unique optimal feasible solution.
 - LP has a nonunique optimal feasible solution (i.e., at least two).

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Solving an LP with Maple.

Farmer Brown:

$$\text{max } 40p + 120w$$

$$\text{s.t. } p + w \leq 100 \quad [\text{land}]$$

$$p + 4w \leq 160 \quad [\text{labor}]$$

$$10p + 20w \leq 1100 \quad [\text{capital}]$$

$$p \geq 0, w \geq 0.$$

In Maple:

restart;

with(Optimization);

$$f := (p, w) \rightarrow 40*p + 120*w;$$

$$\text{Constraints} := [p + w \leq 100, p + 4w \leq 160, \\ 10p + 20w \leq 1100];$$

LPSolve(f(p, w), constraints, 'maximize',
assume = nonnegative);

Output: [5400., [p=60., w=25.]]