# Department of Mathematics Carnegie Mellon University

21-301 Combinatorics, Fall 2005: Test 4

Name:



## Q1: (33pts)

Consider the following general game involving one pile of chips . There is a finite set of positive integers S and each move involves choosing  $s \in S$  and removing this number of chips from the pile. The game ends when there are no moves possible. Show that the SPRAGUE-GRUNDY function  $g$  satisfies  $g(n) \leq |S|$  for all  $n \geq 0$ .

**Solution** First observe that if  $A$  is a set of non-negative integers, then  $mex(A) \leq |A|$ . This is because either  $A = \{0, 1, 2, \ldots, |A| - 1\}$ , in which case  $mex(A) = |A|$  or there exists  $i < |A|$  such that  $i \notin A$ , in which case  $mex(A) \leq i.$ 

Now  $g(n) = mex(A)$  where  $A = \{g(n-s); s \in S\}$  and  $|A| \leq |S|$ .

## Q2: (33pts)

Consider the following game involving two piles of chips. A move consists of removing one pile and splitting the remaining pile into two non-empty piles. There is a unique terminal position in which both piles have one chip. Suppose that the two piles have  $m, n$  chips respectively. Here are the  $N, P$ positions for  $1 \leq m, n \leq 6$ .



When in general is  $m, n$  a  $P$  position. Prove your claim.

**Solution**  $m, n$  is a P-position iff  $m, n$  are both odd. We can prove this by induction on  $m + n$ . The base case is  $m + n = 2$  i.e.  $m = n = 1$ .

If  $m, n$  are both odd then any move will involve splitting one of the numbers into the sum of an odd and an even position. This is an N-position, by induction.

If m is even then we can remove the n and split m into the sum of two odd numbers, a P-position, by induction.

Thus our partition into  $N, P$  has the requisite properties.

### Q3: (34pts)

Consider the following game involving one pile of chips. A move consists of removing  $2^k$  chips where  $k = 0, 1, 2, 3, \ldots$  The first few values of the SPRAGUE-GRUNDY function  $g$  are given in the following table:



What is  $q(n)$  in general? Prove your claim by induction.

**Solution**  $g(n) = n \mod 3$ . It is clearly true for  $n \leq 16$ .

Observe that  $2^k \mod 3 \neq 0$  and so  $n \neq n - 2^k \mod n$  for any  $k \geq 0$ . Since n mod  $3 \neq n-1$  mod  $3 \neq n-2$  mod 3 we see by the induction hypothesis that

$$
g(n) = \max(g(n-1), g(n-2), g(n-4), \ldots) = \max(g(n-1), g(n-2)).
$$

 $n = 3k$ :  $q(n) = mex(3k - 1 \mod 3, 3k - 2 \mod 3) = mex(2, 1) = 0$  $n = 3k + 1: g(n) = mex(3k \mod 3, 3k - 1 \mod 3) = mex(0, 2) = 1$  $n = 3k + 2$ :  $g(n) = mex(3k + 1 \mod 3, 3k \mod 3) = mex(1, 0) = 2$ 

#### Q4: (33pts)

An  $m \times n$  0,1 matrix A and a  $p \times q$  0,1 matrix B are compatible if  $A(i, j) =$  $B(i,j)$  for  $1 \leq i \leq \min\{m,p\}$  and  $1 \leq j \leq \min\{n,q\}.$ 

Suppose that  $A_i$  is an  $m_i \times n_i$  0, 1 matrix for  $i = 1, 2, ..., N$  and that there do **not** exist i, j such that  $A_i$  is compatible with  $A_j$ . Prove that

$$
\sum_{i=1}^{N} \frac{1}{2^{m_i n_i}} \le 1.
$$

**Solution** Let  $P = \max m_i$  and  $Q = \max n_i$ . Let X be a random  $P \times Q$  0,1 matrix and let  $\mathcal{E}_i$  be the event that  $X(p,q) = A_i(p,q)$  for  $1 \leq p \leq m_i, 1 \leq$  $q \leq n_i$ . Then the  $\mathcal{E}_i$  are disjoint events, since if  $\mathcal{E}_i$  and  $\mathcal{E}_j$  both occur, then  $A_i$  and  $A_j$  must be a compatible pair. Thus

$$
1 \geq \sum_{i=1}^N \mathbf{Pr}(\mathcal{E}_i) = \sum_{i=1}^N \frac{1}{2^{m_i n_i}}.
$$