## POLYNOMIAL-TIME PERFECT MATCHINGS IN DENSE HYPERGRAPHS

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A well-known theorem of Rödl, Ruciński and Szemerédi states that any k-graph H on n vertices with minimum codegree  $\delta(H) \geq n/2 + C$  contains a perfect matching. Indeed, for large n the theorem establishes precisely the best-possible value of C for which this statement holds. We therefore cannot be sure that a k-graph H satisfying a weaker minimum degree condition will contain a perfect matching. However, we might hope to prove that such a k-graph either contains a perfect matching or has a given extremal structure for which there is no perfect matching. That is, we would like to characterise those k-graphs which satisfy some weaker minimum degree condition but have no perfect matching.

In this talk I will present such a characterisation for k-graphs H with  $\delta(H) \geq n/k + o(n)$ . For k=3 this characterisation has a simple formulation: for any 3-graph H with  $\delta(H) \geq n/3 + o(n)$ , either H contains a perfect matching or there exists some  $A \subseteq V(H)$  such that |A| is odd but  $|e \cap A|$  is even for any  $e \in E(H)$ . Unfortunately, the naive generalisation of this result for  $k \geq 4$  is false; our characterisation for these values of k is more complicated.

I will also outline a polynomial-time algorithm which tests for this characterisation. As a consequence, we can determine in polynomial time whether or not a k-graph H on n vertices with  $\delta(H) \geq n/k + o(n)$  contains a perfect matching. Furthermore, by derandomising a relatively straightforward random algorithm, we can repeatedly use this testing algorithm to find a perfect matching in such an H in polynomial time (if one exists).

Let  $\mathrm{PM}(k,\delta)$  denote the decision problem of determining whether or not a k-graph H on n vertices with  $\delta(H) \geq \delta n$  contains a perfect matching. So the results described above imply that  $\mathrm{PM}(k,\delta)$  is in P for any  $\delta > 1/k$ . This essentially answers a problem of Karpiński, Ruciński and Szymańska, who had previously shown the existence of  $\varepsilon$  such that  $\mathrm{PM}(k,1/2-\varepsilon)$  is in P. Indeed, Szymańska gave an elegant reduction proving that  $\mathrm{PM}(k,\delta)$  is NP-complete for any  $\delta < 1/k$ , so the minimum codegree threshold at which the perfect matching problem becomes tractable is asymptotically n/k.

This is joint work with Peter Keevash and Fiachra Knox. (For simplicity, the condition  $k \mid n$  has been omitted throughout this abstract; this is a necessary condition for a k-graph to contain a perfect matching.)